

Code Generation for Data Processing

Lecture 5: Analyses and Transformations

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Program Transformation: Motivation

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 - ▶ What parts are actually used? How to find out?
- ▶ Transformation to “better” code must be done *somewhere*

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- ▶ Transformation to “better” code must be done *somewhere*

- ▶ Optimization is a misnomer: we don't know whether it improves code!
 - ▶ Many transformations are driven by heuristics
- ▶ Many types of optimizations are well-known¹¹

¹¹FE Allen and J Cocke. *A catalogue of optimizing transformations*. 1971. .

Dead Block Elimination

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Dead Block Elimination

- ▶ CFG not necessarily connected
- ▶ E.g., consequence of optimization
 - ▶ Conditional branch → unconditional branch
- ▶ Removing dead blocks is trivial
 1. DFS traversal of CFG from entry, mark visited blocks
 2. Remove unmarked blocks

Optimization Example 1

```
define i32 @fac(i32 %0) {  
    br label %for.header  
for.header: ; preds = %for.body, %1  
    %a = phi i32 [ 1, %1 ], [ %a.new, %for.body ]  
    %b = phi i32 [ 0, %1 ], [ %b.new, %for.body ]  
    %i = phi i32 [ 0, %1 ], [ %i.new, %for.body ]  
    %cond = icmp sle i32 %i, %0  
    br i1 %cond, label %for.body, label %exit  
for.body: ; preds = %for.header  
    %a.new = mul i32 %a, %i  
    %b.new = add i32 %b, %i  
    %i.new = add i32 %i, 1  
    br label %for.header  
exit: ; preds = %for.header  
    %absum = add i32 %a, %b  
    ret i32 %a  
}
```

Simple Dead Code Elimination (DCE)

- ▶ Look for trivially dead instructions
 - ▶ No users or side-effects
 - ▶ Calls *might* be removed
- 1. Add all instructions to work queue
- 2. While work queue not empty:
 - 2.1 Check for deadness (zero users, no side-effects)
 - 2.2 If dead, remove and add all operands to work queue

Warning: Don't implement it this naively, this is inefficient

Applying Simple DCE

```
define i32 @fac(i32 %0) {
eff.: cf    br label %for.header
for.header: ; preds = %for.body, %1
users: 2    %a = phi i32 [ 1, %1 ], [ %a.new, %for.body ]
users: 2    %b = phi i32 [ 0, %1 ], [ %b.new, %for.body ]
users: 4    %i = phi i32 [ 0, %1 ], [ %i.new, %for.body ]
users: 1    %cond = icmp sle i32 %i, %0
eff.: cf    br i1 %cond, label %for.body, label %exit
for.body:  ; preds = %for.header
users: 1    %a.new = mul i32 %a, %i
users: 1    %b.new = add i32 %b, %i
users: 1    %i.new = add i32 %i, 1
eff.: cf    br label %for.header
exit:      ; preds = %for.header
users: 0   %absum = add i32 %a, %b
eff.: cf    ret i32 %a
}
```

Applying Simple DCE

```
define i32 @fac(i32 %0) {
eff.: cf    br label %for.header
for.header: ; preds = %for.body, %1
users: 1    %a = phi i32 [ 1, %1 ], [ %a.new, %for.body ]
users: 1    %b = phi i32 [ 0, %1 ], [ %b.new, %for.body ]
users: 4    %i = phi i32 [ 0, %1 ], [ %i.new, %for.body ]
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eff.: cf    ret i32 %a
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```

Dead Code Elimination

- ▶ Problem: unused value cycles

Dead Code Elimination

- ▶ Problem: unused value cycles
 - ▶ Idea: find “value sinks” and mark all needed values as live
 - ▶ Sink: instruction with side effects (e.g., store, control flow)
1. Only mark instrs. with side effects as live
 2. Populate work list with newly added live instrs.
 3. While work list not empty:
 - 3.1 Mark dead operand instructions as live and add to work list
 4. Remove instructions not marked as live

Applying Liveness-based DCE

```
define i32 @fac(i32 %0) {
  br1 label %for.header
for.header: ; preds = %for.body, %1
  %a = phi i32 [ 1, %1 ], [ %a.new, %for.body ]
  %b = phi i32 [ 0, %1 ], [ %b.new, %for.body ]
  %i = phi i32 [ 0, %1 ], [ %i.new, %for.body ]
  %cond = icmp sle i32 %i, %0
  br2 i1 %cond, label %for.body, label %exit
for.body: ; preds = %for.header
  %a.new = mul i32 %a, %i
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  %i.new = add i32 %i, 1
  br3 label %for.header
exit: ; preds = %for.header
  %absum = add i32 %a, %b
  ret i32 %a
}
```

Work list (stack)

Applying Liveness-based DCE

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define i32 @fac(i32 %0) {  
live   br1 label %for.header  
for.header: ; preds = %for.body, %1  
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live   br3 label %for.header  
exit: ; preds = %for.header  
    %absum = add i32 %a, %b  
live   ret i32 %a  
}
```

Work list (stack)

br₁

br₂

br₃

ret

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```
define i32 @fac(i32 %0) {  
live  br1 label %for.header  
for.header: ; preds = %for.body, %1  
live  %a = phi i32 [ 1, %1 ], [ %a.new, %for.body ]  
      %b = phi i32 [ 0, %1 ], [ %b.new, %for.body ]  
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}
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br₁

br₂

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%a

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}
```

Work list (stack)

```
br1  
br2  
br3  
%a.new
```

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for.header: ; preds = %for.body, %1  
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}
```

Work list (stack)

br₁

br₂

br₃

%i

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```
define i32 @fac(i32 %0) {  
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}
```

Work list (stack)

```
br1  
br2  
br3  
%i.new
```

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br₁

br₂

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Work list (stack)

br₁
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Work list (stack)

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Work list (stack)

Liveness-based DCE: Work List Implementation

- ▶ What operations are performed on a work list?
- ▶ How to implement an efficient work list?

Liveness-based DCE: Work List Implementation

- ▶ What operations are performed on a work list?
 - ▶ Insert instruction
 - ▶ Remove any instruction
 - ▶ Test whether instruction is contained
 - ▶ Get and remove next instruction to handle

- ▶ How to implement an efficient work list?

Optimization Example 2

```
define i32 @foo(i32 %0, ptr %1, ptr %2) {  
    %4 = zext i32 %0 to i64  
    %5 = getelementptr inbounds i32, ptr %1, i64 %4  
    %6 = load i32, ptr %5, align 4  
    %7 = zext i32 %0 to i64  
    %8 = getelementptr inbounds i32, ptr %2, i64 %7  
    %9 = load i32, ptr %8, align 4  
    %10 = add nsw i32 %6, %9  
    ret i32 %10  
}
```

Common Subexpression Elimination (CSE) – Attempt 1

- ▶ Idea: find/eliminate redundant computation of same value

Common Subexpression Elimination (CSE) – Attempt 1

- ▶ Idea: find/eliminate redundant computation of same value
- ▶ Keep track of previously seen values in hash map
- ▶ Iterate over all instructions
 - ▶ If found in map, remove and replace references
 - ▶ Otherwise add to map
- ▶ Easy, right?

CSE Attempt 1 – Example 1

```
define i32 @foo(i32 %0, ptr %1, ptr %2) {  
  %4 = zext i32 %0 to i64  
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dup %4  %7 = zext i32 %0 to i64  
        %8 = getelementptr inbounds i32, ptr %2, i64 %7  
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        %10 = add nsw i32 %6, %9  
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→ ht      %5 = getelementptr inbounds i32, ptr %1, i64 %4  
→ ht      %6 = load i32, ptr %5, align 4  
dup %4    %7 = zext i32 %0 to i64  
→ ht      %8 = getelementptr inbounds i32, ptr %2, i64 %4  
→ ht      %9 = load i32, ptr %8, align 4  
→ ht      %10 = add nsw i32 %6, %9  
→ ht      ret i32 %10  
}
```

- ▶ Obsolete instr. can be killed immediately, or in a later DCE

CSE Attempt 1 – Example 2

```
define i32 @square(i32 %a, i32 %b) {
entry:
  %cmp = icmp slt i32 %a, %b
  br i1 %cmp, label %if.then, label %if.end
if.then: ; preds = %entry
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```

Instruction does not dominate all uses!

error: input module is broken!

World Domination

Domination

- ▶ Remember: CFG $G = (N, E, s)$ with digraph (N, E) and entry $s \in N$
- ▶ Dominate: $d \text{ dom } n$ iff every path from s to n contains d
 - ▶ Dominators of n : $DOM(n) = \{d \mid d \text{ dom } n\}$
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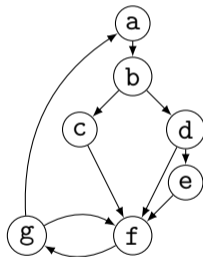
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- \Rightarrow All values not from dominators **not** usable

Dominator Tree

- ▶ Tree of immediate dominators
- ▶ Allows to iterate over blocks in pre-order/post-order
- ▶ Answer $a \text{ sdom } b$ quickly

Control Flow
Graph

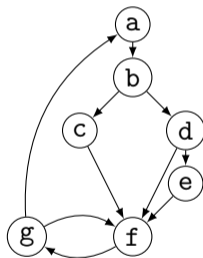


Dominator Tree

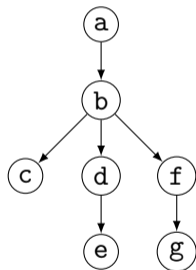
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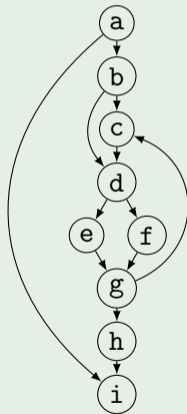
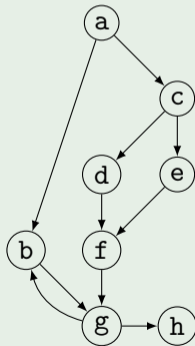


Dominator Tree




Dominator Tree – Example

Construct the dominator tree for the following CFGs (entry at a):




Dominator Tree: Construction


- ▶ Naive: inefficient (but reasonably simple)¹²
 - ▶ For each block: find a path from the root – superset of dominators
 - ▶ Remove last block on path and check for alternative path
 - ▶ If no alternative path exists, last block is idom

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- ▶ Lengauer–Tarjan: more efficient methods¹³
 - ▶ Simple method in $\mathcal{O}(m \log n)$; sophisticated method in $\mathcal{O}(m \cdot \alpha(m, n))$
($\alpha(m, n)$ is the inverse Ackermann function, grows *extremely* slowly)
 - ▶ Used in some compilers¹⁴
- ▶ Semi-NCA: $\mathcal{O}(n^2)$, but lower constant factors¹⁵

¹²ES Lowry and CW Medlock. “Object code optimization”. In: *CACM* 12.1 (1969), pp. 13–22. 

¹³T Lengauer and RE Tarjan. “A fast algorithm for finding dominators in a flowgraph”. In: *TOPLAS* 1.1 (1979), pp. 121–141. 

¹⁴Example: <https://github.com/WebKit/WebKit/blob/aabfacb/Source/WTF/wtf/Dominators.h>

¹⁵L Georgiadis. “Linear-Time Algorithms for Dominators and Related Problems”. PhD thesis. Princeton University, Nov. 2005

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- ▶ Check whether a sdom b ¹⁷
 - ▶ $a.preNum < b.preNum \wedge a.postNum > b.postNum$
 - ▶ After updates, numbers might be invalid: recompute or walk tree

¹⁷PF Dietz. "Maintaining order in a linked list". In: *STOC*. 1982, pp. 122–127. 

Dominator Tree: Implementation

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- ▶ Problem: dominance of unreachable blocks ill-defined \rightsquigarrow special handling

CSE Attempt 2

- ▶ Option 1:
 - ▶ For identical instructions, store all
 - ▶ Add dominance check before replacing
 - ▶ Visit nodes in reverse post-order (i.e., topological order)

- ▶ Option 2:¹⁸
 - ▶ Do a DFS over dominator tree
 - ▶ Use scoped hashmap to track available values

Does this work?

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Does this work? Yes.

CSE: Hashing an Instruction (and Beyond)

- ▶ Needs hash function *and* “relaxed” equality
- ▶ Idea: combine opcode and operands/constants into hash value
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- ▶ Needs hash function *and* “relaxed” equality
- ▶ Idea: combine opcode and operands/constants into hash value
 - ▶ Use pointer or index for instruction result operands
- ▶ Canonicalize commutative operations
 - ▶ Order operands deterministically, e.g., by address
- ▶ Identities: $a+(b+c)$ vs. $(a+b)+c$

Global Value Numbering – or: advanced CSE

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- ▶ Hash-based approach only catches trivially removable duplicates
- ▶ Alternative: partition values into *congruence classes*
 - ▶ Congruent values are guaranteed to always have the same value
- ▶ Optimistic approach: values are congruent unless proven otherwise
- ▶ Pessimistic approach: values are not congruent unless proven
- ▶ Combinable with: reassociation, DCE, constant folding
- ▶ Rather complex, but can be highly beneficial¹⁹

¹⁹K Gargi. "A sparse algorithm for predicated global value numbering". In: *PLDI. 2002*, pp. 45–56.

Simple Transformations: Inlining

- ▶ Estimate whether inlining is beneficial
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- ▶ Replace returns with branches and ϕ -node to/at continuation point
- ▶ Move `alloca` to beginning or save stack pointer
 - ▶ Prevent unbounded stack growth in loops
 - ▶ LLVM provides `stacksave/stackrestore` intrinsics
- ▶ Exceptions may need special treatment

Simple Transformations: Mem2Reg and SROA

- ▶ Mem2reg: promote `alloca` to SSA values/`phis`
 - ▶ Condition: only `load/store`, no address taken
 - ▶ Essentially just SSA construction
 - ▶ Not run in default pipeline, subsumed by SROA

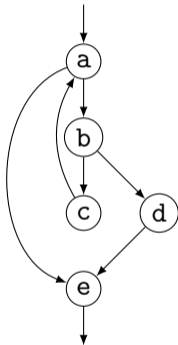
- ▶ SROA: scalar replacement of aggregate
 - ▶ Separate structure fields into separate variables
 - ▶ Also promote them to SSA

What is a Loop?

```
void func() {  
    while (a()) {  
        if (b()) {  
            d();  
            break;  
        }  
        c();  
    }  
    e();  
}
```

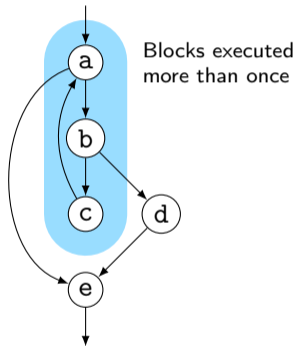
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- ▶ Loops in source code \neq loops in CFG
- ▶ d is *not* part of loop: executed at most once

~> Need algorithm to find loops in CFG

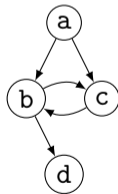
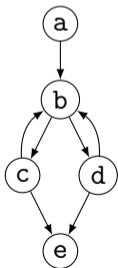
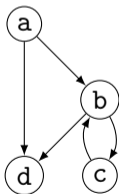
Loops

- ▶ Loop: maximal SCC L with at least one internal edge²⁰
(strongly connected component (SCC): all blocks reachable from each other)
 - ▶ Entry: block with an edge from outside of L
 - ▶ Header h : first entry found (might be ambiguous)
- ▶ Loop nested in L : loop in subgraph $L \setminus \{h\}$

²⁰P Havlak. "Nesting of reducible and irreducible loops". In: *TOPLAS* 19.4 (1997), pp. 557–567. .

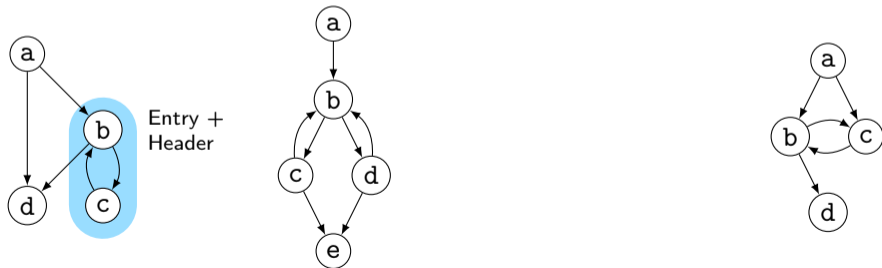
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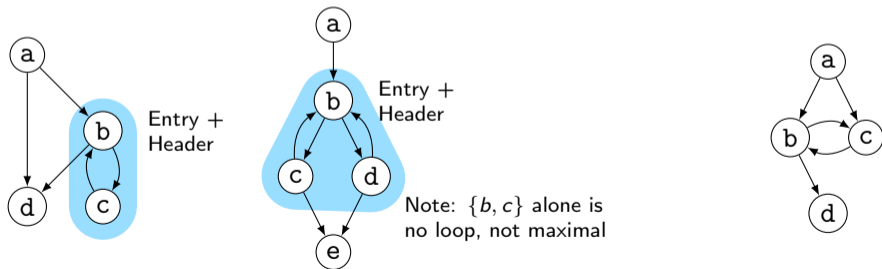
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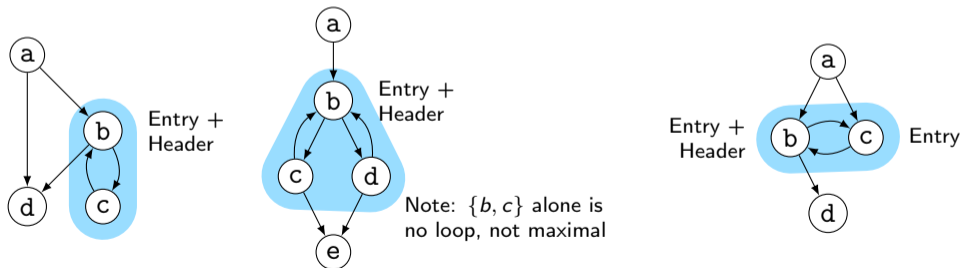
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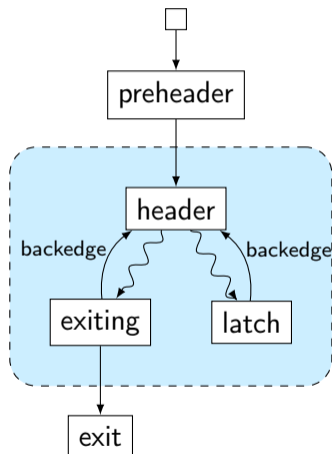
Natural Loops

- ▶ Natural Loop: loop with single entry
 - ⇒ Header is unique
 - ⇒ Header dominates all block
 - ⇒ Loop is reducible
- ▶ Backedge: edge from block to header
- ▶ Predecessor: block with edge into loop
- ▶ Preheader: unique predecessor

Formal Definition


Loop L is reducible iff $\exists h \in L . \forall n \in L . h \text{ dom } n$


CFG is reducible iff all loops are reducible



Finding Natural Loops

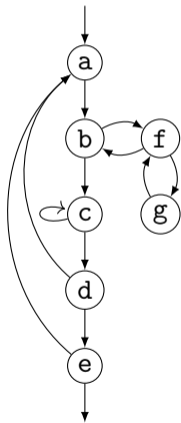
- ▶ Modified version²¹ of Tarjan's algorithm²²
- ▶ Iterate over dominator tree in post order
- ▶ Each block: find predecessors dominated by the block
 - ▶ None \rightsquigarrow no loop header, continue
 - ▶ Any \rightsquigarrow loop header, these edges *must* be backedges
- ▶ Walk through predecessors until reaching header again
 - ▶ All blocks on the way must be part of the loop body
 - ▶ Might encounter nested loops, update loop parent

²¹G Ramalingam. "Identifying loops in almost linear time". In: *TOPLAS* 21.2 (1999), pp. 175–188. .

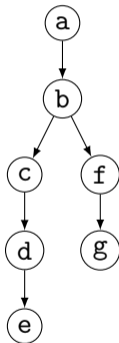
²²R Tarjan. "Testing flow graph reducibility". In: *STOC*. 1973, pp. 96–107. .

Finding Natural Loops: Example

Control Flow Graph



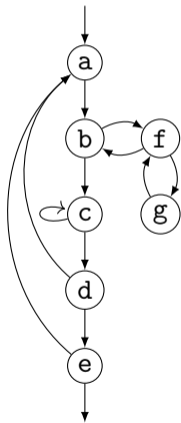
Dominator Tree



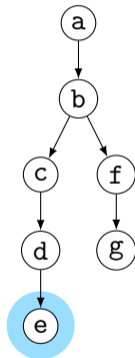
Loop Info

Finding Natural Loops: Example

Control Flow Graph



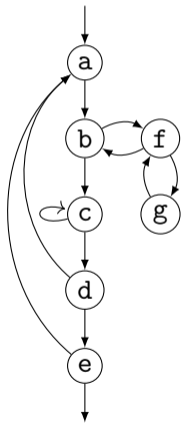
Dominator Tree



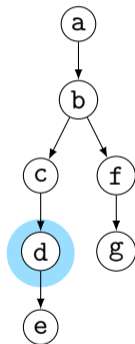
Loop Info

Finding Natural Loops: Example

Control Flow Graph



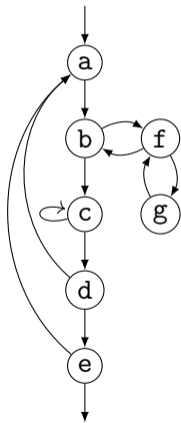
Dominator Tree



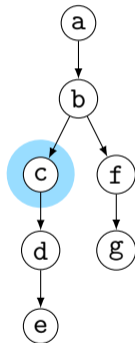
Loop Info

Finding Natural Loops: Example

Control Flow Graph



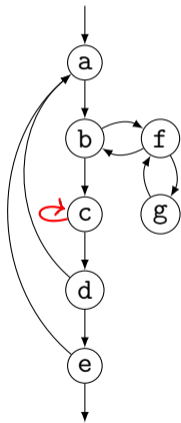
Dominator Tree



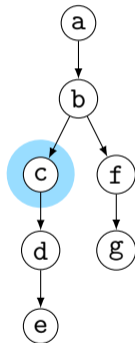
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Finding Natural Loops: Example

Control Flow Graph



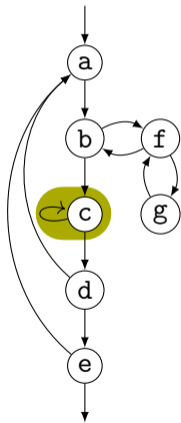
Dominator Tree



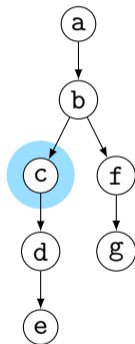
Loop Info

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

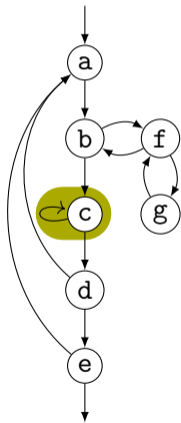


Loop Info

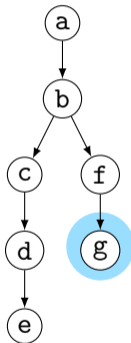
Loop **A**: {c}
header: c; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

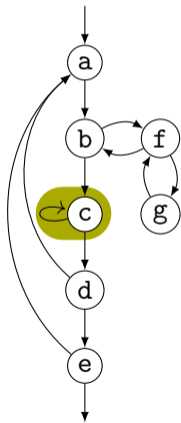


Loop Info

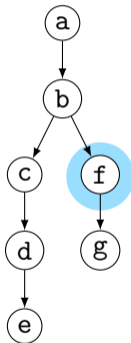
Loop **A**: {c}
header: c; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

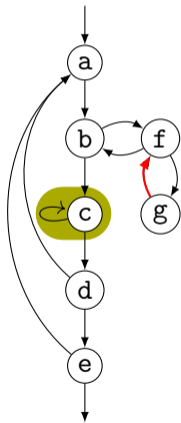


Loop Info

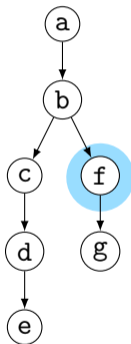
Loop **A**: {c}
header: c; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

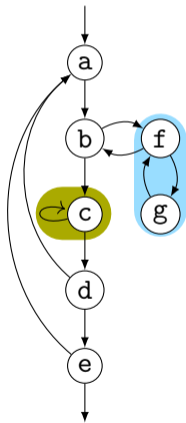


Loop Info

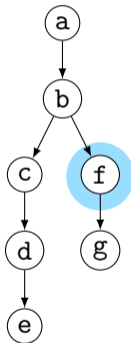
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Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

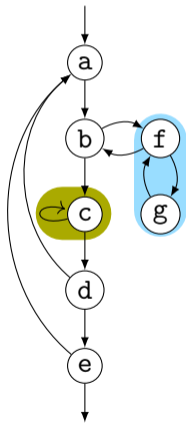


Loop Info

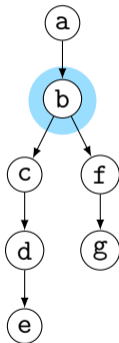
Loop **A**: {c}
header: c; parent: NULL
Loop **B**: {f,g}
header: f; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree



Loop Info

Loop **A**: {c}

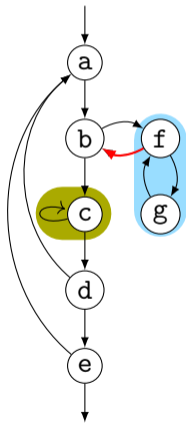
header: c; parent: NULL

Loop **B**: {f,g}

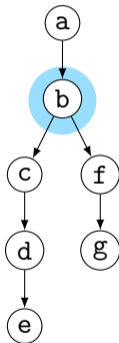
header: f; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree



Loop Info

Loop **A**: {c}

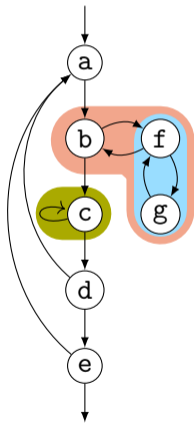
header: c; parent: NULL

Loop **B**: {f,g}

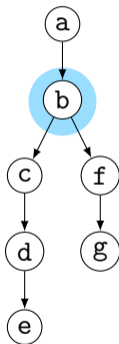
header: f; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

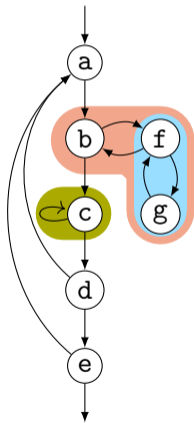


Loop Info

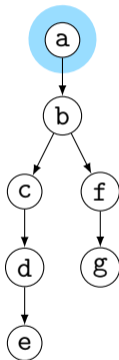
Loop **A**: {c}
header: c; parent: NULL
Loop **B**: {f,g}
header: f; parent: C
Loop **C**: {b,f,g}
header: b; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

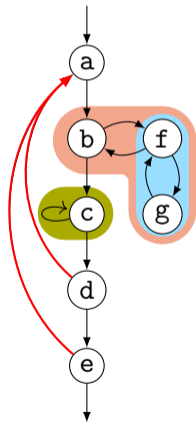


Loop Info

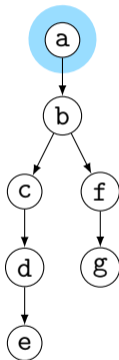
Loop **A**: {c}
header: c; parent: NULL
Loop **B**: {f, g}
header: f; parent: C
Loop **C**: {b, f, g}
header: b; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

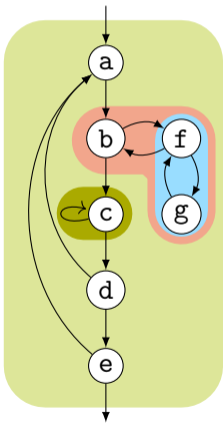


Loop Info

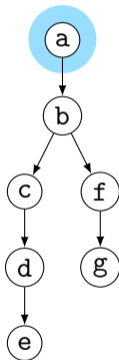
- Loop **A**: {c}
header: c; parent: NULL
- Loop **B**: {f,g}
header: f; parent: C
- Loop **C**: {b,f,g}
header: b; parent: NULL

Finding Natural Loops: Example

Control Flow Graph



Dominator Tree

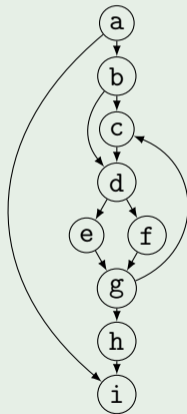
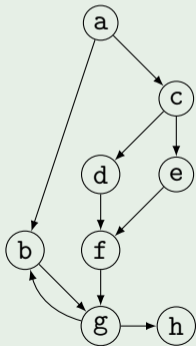


Loop Info

- Loop **A**: {c}
header: c; parent: D
- Loop **B**: {f,g}
header: f; parent: C
- Loop **C**: {b,f,g}
header: b; parent: D
- Loop **D**: {a,b,c,d,e,f,g}
header: a; parent: NULL

Loop Analysis – Example

Apply the previous algorithm to find loops in the following CFGs (entry at a):



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 - ▶ If not, move to preheader (create one, if not existent)
 - ▶ Otherwise, add inst. to set of values defined inside loop

²³<https://github.com/bytecodealliance/wasmtime/blob/bd6fe11/cranelift/codegen/src/licm.rs>

Loop Invariant Code Motion (LICM)

- ▶ Analyze loops, iterate over loop tree in post-order
 - ▶ I.e., visit inner loops first
- ↑ Hoist:²³ iterate over blocks of loop in reverse post-order
 - ▶ For each movable inst., check for loop-defined operands
 - ▶ If not, move to preheader (create one, if not existent)
 - ▶ Otherwise, add inst. to set of values defined inside loop
- ↓ Sink: Iterate over blocks of loop in post-order
 - ▶ For each movable inst., check for users inside loop
 - ▶ If none, move to unique exit (if existent)

²³<https://github.com/bytecodealliance/wasmtime/blob/bd6fe11/craneflift/codegen/src/licm.rs>

Transformations and Analyses in LLVM: Passes

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- ▶ Transformation pass: takes input IR and returns preserved analyses
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- ▶ Transformation pass: takes input IR and returns preserved analyses
 - ▶ Can use analyses, which are re-run when outdated
- ▶ Pass manager executes passes on same granularity
 - ▶ Otherwise, use adaptor: `createFunctionToLoopPassAdaptor`
(and preferably combine multiple smaller passes into a separate pass manager)

Using LLVM (New) Pass Manager

```
void optimize(llvm::Function* fn) {
    llvm::PassBuilder pb;
    llvm::LoopAnalysisManager lam{};
    llvm::FunctionAnalysisManager fam{};
    llvm::CGSCCAnalysisManager cgam{};
    llvm::ModuleAnalysisManager mam{};
    pb.registerModuleAnalyses(mam);
    pb.registerCGSCCAnalyses(cgam);
    pb.registerFunctionAnalyses(fam);
    pb.registerLoopAnalyses(lam);
    pb.crossRegisterProxies(lam, fam, cgam, mam);

    llvm::FunctionPassManager fpm{};
    fpm.addPass(llvm::DCEPass());
    fpm.addPass(llvm::createFunctionToLoopPassAdaptor(llvm::LoopRotatePass()));
    fpm.run(*fn, fam);
}
```

Writing a Pass for LLVM's New PM – Part 1

```
#include "llvm/IR/PassManager.h"
#include "llvm/Passes/PassBuilder.h"
#include "llvm/Passes/PassPlugin.h"

class TestPass : public llvm::PassInfoMixin<TestPass> {
public:
    llvm::PreservedAnalyses run(llvm::Function &F,
                               llvm::FunctionAnalysisManager &AM) {
        // Do some magic
        llvm::DominatorTree *DT = &AM.getResult<llvm::DominatorTreeAnalysis>(F);
        // ...
        llvm::errs() << F.getName() << "\n";
        return llvm::PreservedAnalyses::all();
    }
};
// ...
```

Writing a Pass for LLVM's New PM – Part 2

```
extern "C" ::llvm::PassPluginLibraryInfo LLVM_ATTRIBUTE_WEAK
llvmGetPassPluginInfo() {
    return { LLVM_PLUGIN_API_VERSION, "TestPass", "v1",
            [] (llvm::PassBuilder &PB) {
                PB.registerPipelineParsingCallback(
                    [] (llvm::StringRef Name, llvm::FunctionPassManager &FPM,
                        llvm::ArrayRef<llvm::PassBuilder::PipelineElement>) {
                        if (Name == "testpass") {
                            FPM.addPass(TestPass());
                            return true;
                        }
                        return false;
                    });
            } });
}
```

```
c++ -shared -o testpass.so testpass.cc -lLLVM -fPIC
```

```
opt -S -load-pass-plugin=$PWD/testpass.so -passes=testpass input.ll
```

Analyses and Transformations – Summary

- ▶ Program Transformation critical for performance improvement
- ▶ Code not necessarily better
- ▶ Analyses are important to drive transformations
 - ▶ Dominator tree, loop detection, value liveness
- ▶ Important optimizations
 - ▶ Dead code elimination, common sub-expression elimination, loop-invariant code motion
- ▶ Compilers often implement transformations as passes
- ▶ Analyses may be invalidated by transformations, needs tracking

Analyses and Transformations – Questions

- ▶ Why is “optimization” a misleading name for a transformation?
- ▶ How to find unused code sections in a function’s CFG?
- ▶ Why is a liveness-based DCE better than a simple, user-based DCE?
- ▶ What is a dominator tree useful for?
- ▶ What is the difference between an irreducible and a natural loop?
- ▶ How to find natural loops in a CFG?
- ▶ How does the algorithm handle irreducible loops?
- ▶ Why is sinking a loop-invariant inst. harder than hoisting?